

Today

Spectral Methods (Part II)

Adjacency Matrix

- Writ's Theorem
- Hoffman Bd

Laplacian

- Drawing Graphs

CSS.205.1

Toolkit in TCS

- Lecture #21

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Spectrum of the Adjacency Matrix

Graph Colouring:

Greedy Algorithm:

1. Order the vertices according to some perm π
2. Colour each vertex (in order accg to π) w/ a colour different from its neighbours that have already been coloured.

$$\chi(G) \leq d_{\max} + 1$$

(chromatic # graph)

π - number k_π s.t



$$k_\pi = \max_u \#\{v \in N(u) \mid \pi(v) < \pi(u)\}$$

- Greedy: $\chi(G) \leq k_\pi + 1$ (if perm π is chosen)

$$\chi(G) \leq \min_{\pi} (k_\pi) + 1$$

Will show that there is a permutation π s.t. $k_\pi \leq \mu_1$ (*)


Theorem: [Wilf's Theorem] $\chi(G) \leq \lfloor \mu_1 \rfloor + 1$ (improvement over diam bound)

Pf: Suffices to show (*).

Last time: $\mu_1 \geq \max_{\emptyset \neq S \subseteq V} d_{ave}(G[S])$

Order vertices in reverse as follows

$\mu_1 \geq d_{ave}$, $\Rightarrow \exists$ vertex v . $\deg(v) \leq \mu_1$

π : 

Now, for remaining vertices.

$G \setminus \{v\}$. $\mu_1 \geq d_{ave}(G \setminus \{v\})$

$\exists v_{n-1}$ s.t. $\deg_{G \setminus \{v\}}(v_{n-1}) \leq \mu_1$

This π satisfies. $k_\pi \leq \mu_1$ \square

Hoffman Bound.

Thm: $\alpha(G) \geq 1 - \frac{\mu_1}{\mu_n} = \frac{\mu_1 - \mu_n}{-\mu_n}$

Today, we will prove the Hoffman bound for d -regular graphs.

Prove a stronger statement.

Thm: $\alpha(G) \leq -\frac{\mu_n}{\mu_1 - \mu_n}$ for d -regular graphs

fractional size of the largest independent set
- $\frac{\text{largest ind set size}}{\# \text{ vertices}}$

[Cor: $\alpha(G) \geq \frac{\mu_1 - \mu_n}{-\mu_n}$]

Pf: G - d -regular

$\Rightarrow \mu_1 = d$ & $\phi_1 = \frac{1}{\sqrt{n}} \mathbb{1}$ ($n = \# \text{ vertices}$)

Let S be any independent set (say of largest size)

$|S| = \alpha n.$

$$f: V \rightarrow \mathbb{R}$$

Indicator fn of the set S .

$$f(v) = \mathbb{1}[v \in S]$$

View $f \in \mathbb{R}^V$

$$- f = \sum_{i=1}^n f_i \varphi_i \quad (\text{where } \varphi_i \text{ are the orthonormal eigen vectors of } A)$$

$$- \langle f, Af \rangle = \sum_{(u,v) \in E} f(u)f(v) = 0 \quad (f \text{ is the indicator of independent set})$$

$$Af = A(\sum f_i \varphi_i) = \sum \mu_i f_i \varphi_i$$

$$\begin{aligned} \langle f, Af \rangle &= \langle \sum f_i \varphi_i, \sum \mu_i f_i \varphi_i \rangle \\ &= \sum_{i,j} f_i \mu_j f_j \langle \varphi_i, \varphi_j \rangle \\ &= \sum_i \mu_i f_i^2 \end{aligned}$$

$$0 = \sum \mu_i f_i^2$$

$$f_i = \langle f, \varphi_i \rangle = \langle f, \frac{1}{\sqrt{n}} \mathbb{1} \rangle = \frac{|S|}{\sqrt{n}} = \alpha \sqrt{n}$$

$$\langle f, f \rangle = \sum_{u \in V} f(u)f(u) = |S| = \alpha \sqrt{n}$$

$$\langle f, f \rangle = \sum f_i^2$$

$$- f_1 = \alpha \sqrt{n}$$

$$- \sum f_i^2 = \alpha n$$

$$\begin{aligned} 0 &= \sum \mu_i f_i^2 = \mu_1 f_1^2 + \sum_{i=2}^n \mu_i f_i^2 \\ &\geq \mu_1 \alpha^2 n + \mu_n \sum_{i=2}^n f_i^2 \\ &= \mu_1 \alpha^2 n + \mu_n \left(\sum_{i=1}^n f_i^2 - f_1^2 \right) \\ &= \mu_1 \alpha^2 n + \mu_n (\alpha n - \alpha^2 n) \end{aligned}$$

Rearranging $\alpha \leq \frac{-\mu_n}{\mu_1 - \mu_n}$ \square

The above proof works if we replace adj matrix A w/ any other matrix B s.t

$$B \mathbb{1} = \mathbb{1} \quad \& \quad B(y_j) = 0 \text{ if } (y_j) \notin E$$

(B - has real eigen spectrum)

$$\alpha(G) \leq \Theta_{\#}(G) \triangleq \min_{\substack{B \text{-symmetric} \\ \lambda_{\min} \\ B(y_j) = 0 \text{ if } (y_j) \notin E \\ B \mathbb{1} = \mathbb{1} \\ B \succcurlyeq \lambda_{\min} I}} \left(\frac{-\lambda_{\min}}{1 - \lambda_{\min}} \right)$$

Lovasz further strengthened.

$$\alpha(G) \leq \alpha_2(G) \equiv \min_{\substack{\text{sym } C \\ \lambda \in \mathbb{R}}} (\lambda)$$

Lovasz Theta
Function.

$$\begin{cases} C_{ij} = 0 & \forall (i,j) \notin E \\ A + C \leq \lambda I \end{cases}$$

Laplacian Matrix & Its Spectrum

$$L_G = D_G - A_G.$$

$$\langle x, L_G x \rangle = \sum_{\{u,v\} \sim E} (x(u) - x(v))^2$$

$$\text{Weighted Graph} \quad \sum_{\{u,v\} \sim E} \omega(u,v) (x(u) - x(v))^2$$

Graph ω / non-negative, L_G is a PSD

$$\text{Eigen values } 0 = \lambda_1 \leq \lambda_2 \leq \dots \leq \lambda_n \in \mathbb{R}$$

$$\text{Eigen vectors } \psi_1 \dots \psi_n \in \mathbb{R}^n$$

Least e-value = Corresponding e-vector

$$\begin{aligned}\lambda_1 &= \min_{x \in \mathbb{R}^n} \frac{\langle x, L_G x \rangle}{\langle x, x \rangle} \\ &= \min_x \frac{\sum_{\{u,v\} \in E} (x(u) - x(v))^2}{\sum_{u \in V} (x(u))^2}\end{aligned}$$

$\lambda_1 = 0$ corresponding e-vector
 $\varphi_1 = \frac{1}{\sqrt{n}} \mathbb{1}$.

If x is an eigenvector of L_G corresponding to eigenvalue 0, then x must be constant on each (connected) component of G .

Multiplicity of 0-eigenvalue = # components of G .

G is connected $\Rightarrow \lambda_2 > 0$.

(λ_2 - measure of how well the graph is connected)

\hookrightarrow Cheeger inequalities.

Hall's drawing of graphs:

One-dim drawing of graphs.

Goal: Plot the vertices on a straight line such that it minimizes the sum of squared edge distances.

Formulation: Find $x: V \rightarrow \mathbb{R}$
st $\sum_{\{u,v\} \in E} (x(u) - x(v))^2$ is minimized.

- Obs: ① Assume

x is normalized i.e.
 $\|x\|^2 = \langle x, x \rangle = 1.$

②. $x = \text{const}$ minimizes.

(not a nice picture)

since all vertices are mapped to same point)

$x \perp \mathbb{1}$ (x is orthogonal to constant vector).

$$- (3) \min_{\substack{x \in \mathbb{R}^V \\ x \perp \mathbb{1} \\ \langle x, x \rangle = 1}} \sum_{\{u, v\} \sim E} (x(u) - x(v))^2 = \lambda_2$$

a corresponding x is ψ_2 - second eigenvector.

$\psi_2: V \rightarrow \mathbb{R}$ is the best one-dim drawing.

What about 2-dim

Find $x: V \rightarrow \mathbb{R}$
 $y: V \rightarrow \mathbb{R}$ st

$$(*) = \sum_{\{u, v\} \sim E} \left[\begin{pmatrix} x(u) \\ y(u) \end{pmatrix} - \begin{pmatrix} x(v) \\ y(v) \end{pmatrix} \right]^2$$

$$= \sum_{\{u, v\} \sim E} \left[(x(u) - x(v))^2 + (y(u) - y(v))^2 \right]$$

For the same reasons as before
 $\|x\|^2 = \|y\|^2 = 1$, $\langle x, \mathbb{1} \rangle = \langle y, \mathbb{1} \rangle = 0$.

Best soln $x = y = \psi_2$

- degenerate 1-dim drawing.

To prevent, $x \perp y$. $\langle x, y \rangle = 0$

This $\left. \begin{array}{l} x = \psi_2 \\ y = \psi_3 \end{array} \right\}$ as candidate vectors

Thm: $x_1, \dots, x_k: V \rightarrow \mathbb{R}$ is a k -dim drawing of the vertices such that

$$- \|x_i\|^2 = 1$$

$$- \langle x_i, x_j \rangle = 0 \quad \forall i \neq j$$

$$- \langle x_i, \mathbb{1} \rangle = 0 \quad \forall i$$

then:

$$\sum_{\{u,v\} \in E} \sum_{i=1}^k (x_i(u) - x_i(v))^2 \geq \sum_{j=2}^{k+1} \lambda_j$$

Spectra of Toy Graphs.

1. Complete Graph K_n

$$A_{K_n}(i,j) = \begin{cases} 1 & \text{if } i \neq j \\ 0 & \text{o.w.} \end{cases}$$

$$L_{K_n} = (\text{Laplacian}), \quad L_{K_n}(i,j) = \begin{cases} -1 & \text{if } i \neq j \\ n-1 & \text{if } i=j \end{cases}$$

$$L_{K_n} \mathbb{1} = \bar{0} \quad (\lambda_1 = 0; \psi_1 = \frac{1}{\sqrt{n}} \mathbb{1})$$

$\varphi \perp \mathbb{1}$

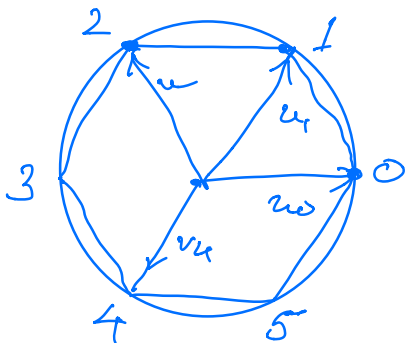
$$\begin{aligned} (L_n \varphi)(v) &= (n+1)\varphi(v) - \sum_{u \neq v} \varphi(u) \\ &= (n-1)\varphi(v) + \varphi(v) \\ &= n\varphi(v). \end{aligned}$$

Hence, any $\varphi \perp \mathbb{1}$, is an e-vector
w/ e-value n .

L_n : Spectrum: e-value $\begin{cases} 0 & \text{w/ mult } 1 \\ n & \text{w/ mult } n-1 \end{cases}$

② Cycle (C_n) / Ring Graph.

$$A_n(u, j) = \begin{cases} 1 & \text{if } |u-j| = 1 \pmod{n} \\ 0 & \text{otherwise.} \end{cases}$$



$$u_0 \dots u_n$$

$$u_i = \left(\cos \frac{2\pi i}{n}, \sin \frac{2\pi i}{n} \right)$$

$$u_{i-1} + u_{i+1} = 2 \cos \frac{2\pi}{n} u_i$$

$$\begin{aligned} L_G &= dI - A_G \\ A_G \varphi &= \lambda \varphi \\ L_G \varphi &= (d - \lambda) \varphi \end{aligned}$$

$$x(i) = \cos \frac{2\pi i}{n} \quad ; \quad i=0, \dots, n-1$$

$$y(i) = \sin \frac{2\pi i}{n} \quad ; \quad i=0, \dots, n-1$$

x, y are e-vectors of A_n w/
e-value. $2 \cos \frac{2\pi}{n}$

$$x_k(i) = \cos \frac{2\pi ki}{n} \quad ; \quad i=0, \dots, n-1$$

$$y_k(i) = \sin \frac{2\pi ki}{n} \quad ; \quad i=0, \dots, n-1$$

$x_k = y_k$ are e-vectors of A_n w/
e-value $2 \cos \frac{2\pi k}{n}$

Eigen-spectrum $1 \leq k \leq n/2$