

Today

Lovasz Theta
Function

CSS.205.1

Toolkit in TCS

- Lecture # 28

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Instructor: Prahladh
Harsha

Recap from last lecture.

Shannon Capacity of a Graph

$G = (V, E)$ - simple graph (no self-loops
& no multiedges)

$$\Sigma(G) = \lim_{k \rightarrow \infty} \sqrt[k]{\alpha(G^{\boxtimes k})} = \sup_{k \rightarrow \infty} \sqrt[k]{\alpha(G^{\boxtimes k})}$$

where strong product $G_1 = (V_1, E_1); G_2 = (V_2, E_2)$

$$G_1 \boxtimes G_2 = (V, E)$$

$$V = V_1 \times V_2$$

$$((u_1, u_2), (v_1, v_2)) \in E \Leftrightarrow (i) (u_1, u_2) \neq (v_1, v_2)$$

$$(ii) \forall i \in [2], u_i = v_i \text{ \& }$$

$$(u_i, v_i) \in E$$

Qn: What is $\Sigma(G)$?

Observations:

- $\alpha(G)$ - super multiplicative w.r.t strong prod
i.e. $\alpha(G \boxtimes H) \geq \alpha(G) \cdot \alpha(H)$ ✓

- Suppose $f: \{\text{graphs}\} \rightarrow \mathbb{R}_{\geq 0}$ s.t

(i) $\alpha(G) \leq f(G)$

(ii) f is sub-multiplicative

i.e. $f(G \boxtimes H) \leq f(G) \cdot f(H)$

then $\Sigma(G) \leq f(G)$.

- Instantiating f with $\bar{\chi}(G)$ (or $\chi(G)$)
 $\Sigma(G) \leq \bar{\chi}(G)$

Hence, $\sqrt{5} \leq \Sigma(G) \leq 3$.

Today: Lovasz's improved UB on $\Sigma(G)$.

Orthogonal Mappings

$i \mapsto v_i$ (unit vectors) | No restriction on ambient dim

$\vec{v} = (v_1, \dots, v_n)$ - collection of unit vectors.

s.t. $\{i, j\} \notin E \Rightarrow v_i \perp v_j$
= $i \neq j$ (orthogonal)

Properties of Orthogonal Mappings

① \bar{u}, \bar{v} - orthogonal mappings to $G \rightarrow H$

\Downarrow

$\bar{u} \otimes \bar{v}$ - orthogonal mapping to $G \otimes H$

$$\bar{u} = (u_1, \dots, u_n) \in \mathbb{R}^n, \quad \bar{v} = (v_1, \dots, v_m) \in \mathbb{R}^m$$

$$\bar{u} \otimes \bar{v} = (\omega_{ij} \mid i \in [n], j \in [m]) \in \mathbb{R}^{nm}$$

$$\omega_{ij} = u_i \otimes v_j \quad (\omega_{ij})$$

$$\omega_{ij}^{(G \otimes H)} = u_i^{(G)} \cdot v_j^{(H)}$$

$$\langle u_i \otimes v_j, u_k \otimes v_l \rangle = \langle u_i, u_k \rangle \cdot \langle v_j, v_l \rangle$$

c - unit vector. (handle).

$$\vartheta(G, \bar{u}, c) = \max_{i \in [n]} \frac{1}{\langle c, u_i \rangle^2}$$

② $\alpha(G) \leq \vartheta(G, \bar{u}, c)$, \forall orthogonal mappings \bar{u} to handle c

Pf: $S \subseteq V$, S is an independent set
 $|S| \leq \alpha(G)$.

$$1 = \langle c, c \rangle \geq \sum_{c \in S} \langle c, u_i \rangle^2 \geq \left(\min_{c \in S} \langle c, u_i \rangle^2 \right) \cdot |S|$$

$$\geq \left(\min_c \langle c, u_i \rangle^2 \right) \cdot |S|$$

$$(3) \quad \vartheta(G \boxtimes H, \bar{u} \otimes \bar{v}, c \otimes d) = \vartheta(G, \bar{u}, c) \cdot \vartheta(H, \bar{v}, d)$$

$$\langle u_i \otimes v_j, c \otimes d \rangle = \langle u_i, c \rangle \cdot \langle v_j, d \rangle$$

— Lovasz' Theta Function:

$$\vartheta(G) = \min_{\bar{u}, c} \vartheta(G, \bar{u}, c)$$

$$(4) \quad \alpha(G) \leq \vartheta(G)$$

$$(5) \quad \vartheta(G \boxtimes H) \leq \vartheta(G) \cdot \vartheta(H)$$

— Corollary: $\Sigma(G) \leq \vartheta(G)$.

— Return to C_5 (5-cycle).

$3 \leq \Sigma(C_5) \leq \vartheta(C_5)$: What is $\vartheta(C_5)$?

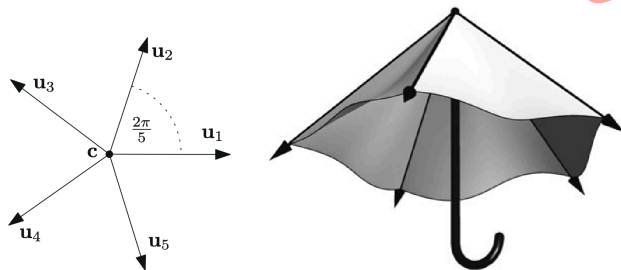


Fig. 3.3 A five-rib umbrella: fully open and viewed from the top (left), and partially folded so that non-adjacent ribs are perpendicular (right)

Candidate Orthogonal Mapping to G_5 .

$$k \in \{0, 1, \dots, 4\}$$

$$u_k = \frac{(\cos \frac{2\pi k}{5}, \sin \frac{2\pi k}{5}, z)}{\sqrt{\cos^2 \frac{2\pi k}{5} + \sin^2 \frac{2\pi k}{5} + z^2}} = \frac{(e^{i2\pi k/5}, z)}{\sqrt{1+z^2}}$$

$$c = (0, 0, 1).$$

Find z st $\langle u_i, u_{i+2} \rangle = 0$

$$\langle u_0, u_2 \rangle = 0$$

$$\langle u_0, u_2 \rangle = \frac{\cos \frac{4\pi}{5} + z^2}{1+z^2}; \text{ Hence } z^2 = -\cos \frac{4\pi}{5}$$

$$= \cos \frac{\pi}{5}$$

$$= \frac{\sqrt{5}+1}{4}$$

$$\vartheta(G_5, \bar{u}, c) = \frac{1}{\langle c, u_0 \rangle^2} = \frac{1+z^2}{z^2} = \frac{(\sqrt{5}+1)/4}{(\sqrt{5}+1)/4} = \sqrt{5}$$

Hence, $\vartheta(G_5) \leq \sqrt{5}$.

Thus, $\Sigma(G_5) = \sqrt{5}$.

□

SDP Formulation for $\vartheta(G)$:

$$\vartheta(G) = \min_{\text{sym } M} \lambda \text{ st } \begin{cases} M_{ij} = 1 \text{ if } i=j \text{ or } i+j \\ M \preceq \lambda I \end{cases}$$

(Pf uses the fact that $N \succeq 0$ (Prop 5)
then $\exists v_1 \dots v_n$ st $N_{ij} = \langle v_i, v_j \rangle$)

Proof will yield an orthonormal mapping
 \Rightarrow handle st

$$\vartheta(G) = \frac{1}{\langle c, v \rangle^2} = \frac{1}{\langle c, v \rangle^2} = \frac{1}{\langle c, v \rangle^2}.$$

Thm [Sandwich Theorem]
 $\alpha(G) \leq \vartheta(G) \leq \bar{\chi}(G) = \chi(\bar{G})$

Pf: Sufficient to prove.

$$\vartheta(\bar{G}) \leq \chi(G).$$

i.e., SDP formulation for $\vartheta(\bar{G})$

$$\vartheta(\bar{G}) = \min_{\text{sym } M} \lambda$$
$$\begin{cases} M_{ij} = 1 & \text{if } i=j \text{ or } \{i,j\} \in E \\ M \preceq \lambda I \end{cases}$$

Suffices to show: Given a k -colouring
of \bar{G} , there is a matrix M satisfying
the above properties $\Rightarrow M \preceq kI$

k -simplex: $e_1 \dots e_k$ - unit vectors.

$$\tilde{e}_i = e_i - \frac{1}{k} \sum_{i=1}^k e_i$$

$$v_i = \frac{\tilde{v}_i}{\|\tilde{v}_i\|}$$

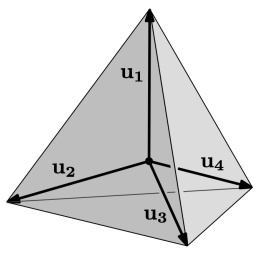
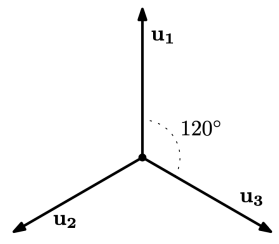


Fig. 3.4 Unit vectors with pairwise scalar products $-1/(k-1)$ for $k=3,4$

$$\langle \tilde{v}_i, \tilde{v}_j \rangle$$

$$\tilde{v}_i = e_i - \frac{1}{k} \sum_{e \in E} e = \left(1 - \frac{1}{k}\right) e_i - \frac{1}{k} \sum_{e \neq e_i} e$$

$$i=j: \|\tilde{v}_i\|^2 = \left(1 - \frac{1}{k}\right)^2 + \frac{(k-1)}{k^2} = \frac{k^2 - 2k + 1 + k - 1}{k^2} = \frac{k-1}{k}$$

$$i \neq j: \langle \tilde{v}_i, \tilde{v}_j \rangle = \frac{-2}{k} \left(1 - \frac{1}{k}\right) + \frac{(k-2)}{k^2}$$

$$= \frac{-2k + 2 + k - 2}{k^2} = -\frac{1}{k}$$

$$\langle v_i, v_j \rangle = \begin{cases} 1 & \text{if } i=j \\ -\frac{1}{k-1} & \text{if } i \neq j \end{cases}$$

$\chi: V \rightarrow \{1, \dots, k\}$ is a k -colouring of \bar{G}

Vector Colouring.

Vertex $i \rightarrow$ Vector $v_i \triangleq v_{\chi(i)}$

Vector Colouring satisfies.

$$(i) \|v_i\| = 1$$

$$(ii) \{i, j\} \in E, \langle v_i, v_j \rangle = -\frac{1}{k-1}$$

Define matrix $N \in \mathbb{R}^{v \times v}$

$$N(i, j) = \langle v_i, v_j \rangle$$

$$\begin{cases} N_{ii} = 1 \\ N_{ij} = -\frac{1}{k-1} \text{ if } \{i, j\} \in E(G) \end{cases}$$

Certainly $N \succeq 0$

Define: $M = kI - (k-1)N$

M satisfies

$$\begin{cases} (i) M \text{ is symmetric} \\ (ii) M_{ii} = k - (k-1) = 1 \\ (ii) M_{ij} = -(k-1) \left(-\frac{1}{k-1} \right) \text{ if } \{i, j\} \in E \\ = 1 \end{cases}$$

$$(iv) M \preceq kI. \quad (k, kI - M \succeq 0)$$

Hence, $\vartheta(\bar{G}) \leq k$.

$$\overline{\vartheta(\bar{G})}; \quad M \rightarrow N = \frac{1}{\lambda-1} (\lambda I - M) \quad \square$$

$\mathcal{V}(G)$ - yields a vector colouring of

$$\subseteq \text{st} \\ \forall i, \|v_i\| = 1$$

$$\forall i \neq j, \langle v_i, v_j \rangle = -\frac{1}{\lambda-1}$$

Further properties of $\mathcal{V}(G)$.

$$\textcircled{1} \quad \mathcal{V}(G) = \max_{\text{sym } B} \text{Tr}(BT) \quad (T \text{ is all } 1\text{'s matrix})$$
$$\left\{ \begin{array}{l} B \succeq 0 \\ B_{ij} = 0 \text{ if } i \neq j \neq i \neq j \\ \text{Tr}(B) = 1 \end{array} \right.$$

$$\textcircled{2} \quad \mathcal{V}(G) = \max_{\substack{\bar{u}, c \\ \text{orth map}}} \sum_{i=1}^n \langle c, u_i \rangle^2$$

$$\textcircled{3} \quad \mathcal{V}(G) = \max_A \left(1 - \frac{\lambda_1(A)}{\lambda_n(A)} \right)$$
$$\left\{ \begin{array}{l} A_{ij} = 0 \text{ if } i \neq j \\ \lambda_1(A) \geq \dots \geq \lambda_n(A) \\ \text{— non-decreasing order} \end{array} \right.$$

Hoffman Bound:

A - adjacency matrix. $\chi(G) \geq 1 - \frac{\mu_1}{\mu_n}$

Pf. $\chi(G) \geq \nu(\bar{G}) \geq 1 - \frac{\mu_1}{\mu_n}$ \square

Sandwich Theorem: $\alpha(G) \leq \nu(G) \leq \bar{\chi}(G)$

Perfect Graphs: G is perfect if
 \forall induced graphs G' , $\alpha(G') = \bar{\chi}(G')$

Weak Perfect Graph Theorem:

G is perfect $\Leftrightarrow \bar{G}$ is perfect