

Today

- High Dimensional  
Expanders

CSS. 413.1

Pseudorandomness

Lecture 27 (2021-12-7)

Instructor: Prahladh  
Harsha.

- High-dimensional expansion

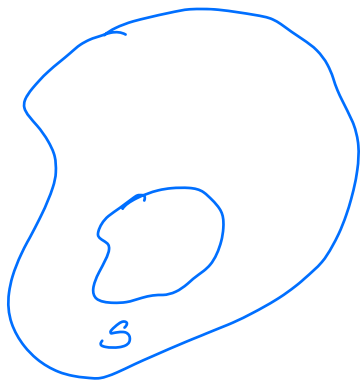
- gen of expansion in graphs  
to hypergraphs

Recap

Expanders Graphs.

- Definition

$G = (V, A)$



(i) Spectral Expansion

$A$ -adj matrix of a  
graph. (normalized).

$$\lambda(A) = \max\{\lambda_2, |\lambda_n|\}$$

$$\lambda \ll 1$$

(ii) Combinatorial Expansion

$$\forall S, |S| \leq k, |N(S)| \geq \alpha |S|$$

(then  $G$  is  $(k, \alpha)$ -expander.

### (iii) Random Walk

- random walk converges to the uniform (stationary) dist fast

-  $A$  - normalized adj matrix

$$J/n = \begin{bmatrix} 1/n & \dots & 1/n \\ \vdots & & \vdots \\ 1/n & & 1/n \end{bmatrix} \quad n = \# \text{ vertices}$$

(random walk independent vertex)

$$\|A - J/n\| \leq \lambda \quad (\lambda \ll 1)$$

then  $A$  is a 'good' expander

(equiv to spectral defn)

Thm: There exists explicit construction of constant degree expanders

→ Is there a similar theory of expansion for hypergraphs?

$$H = (V, F)$$

$\hookrightarrow V$  - vertices  
 $F$  - set of hyperedges

$F \subseteq \binom{V}{k} \quad (k \geq 3)$

}  $k$ -uniform

In case of graphs, the gold-standard - complete graph.

Similarly, for  $k$ -uniform hypergraphs gold standard - complete  $k$ -uniform hypergraph

$(V, \binom{V}{k})$

$k$ -uniform hypergraphs (all edges have same arity)

Multiple Generalization of expansion to hypergraph

- Most generalizations are not equivalent

Today's lecture (high-dimensional expanders)

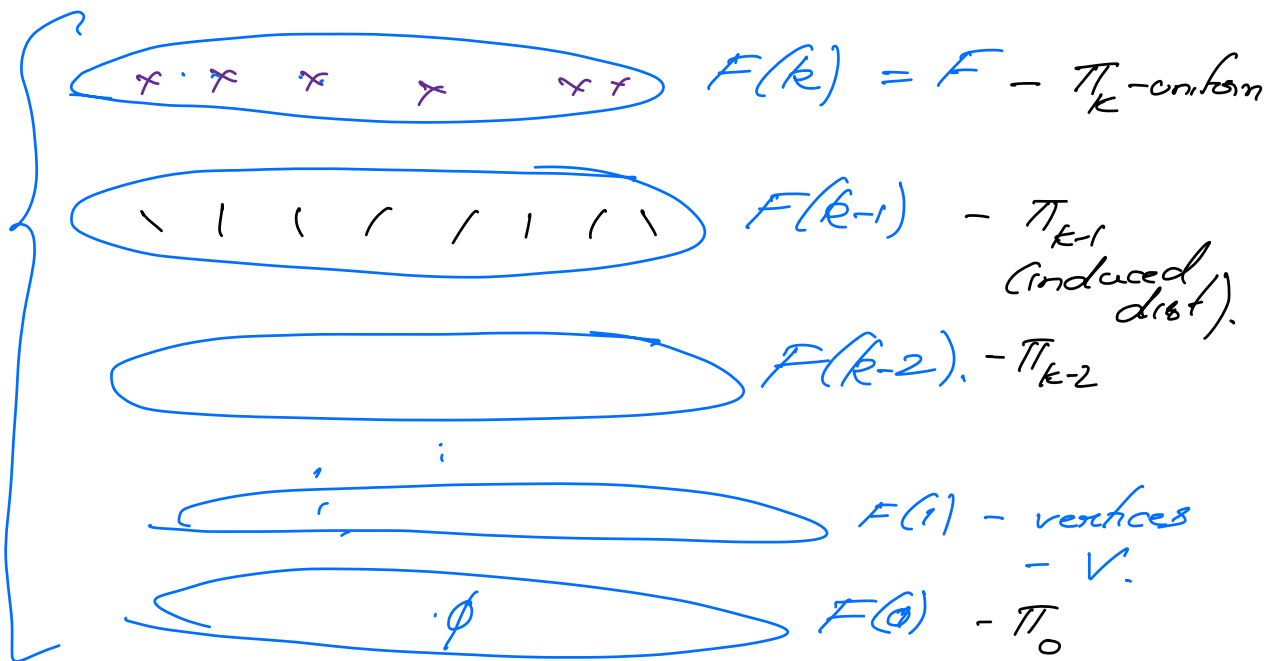
- a generalization based on random walks / spectral.

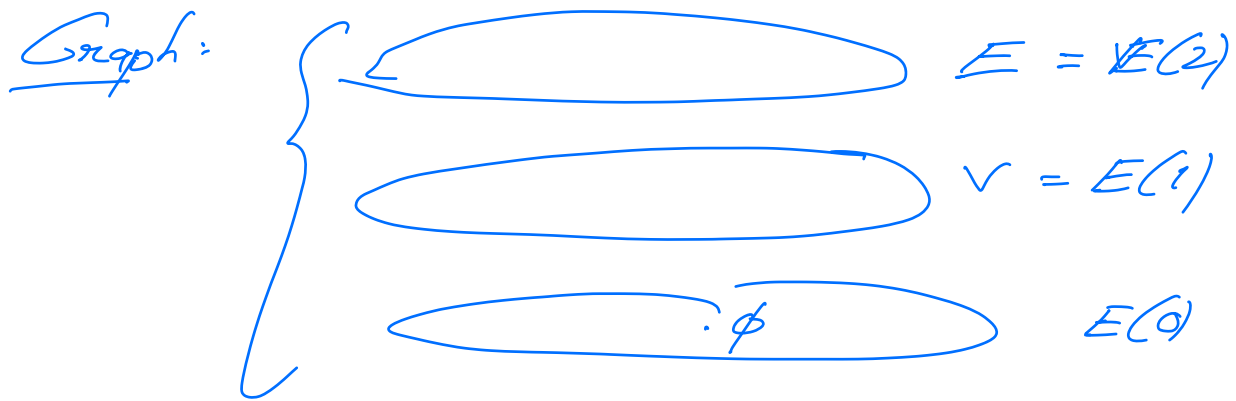
Simplicial viewpoint of hypergraph

$$(V, F) \quad F \subseteq \binom{V}{k}$$

$F = F(k)$  - set of hyperedges of size  $k$ ,  
 $\forall 0 \leq i < k$ ,

$$F(i) = \{S \subseteq \binom{V}{i} \mid \exists T \in F, S \subseteq T\}$$

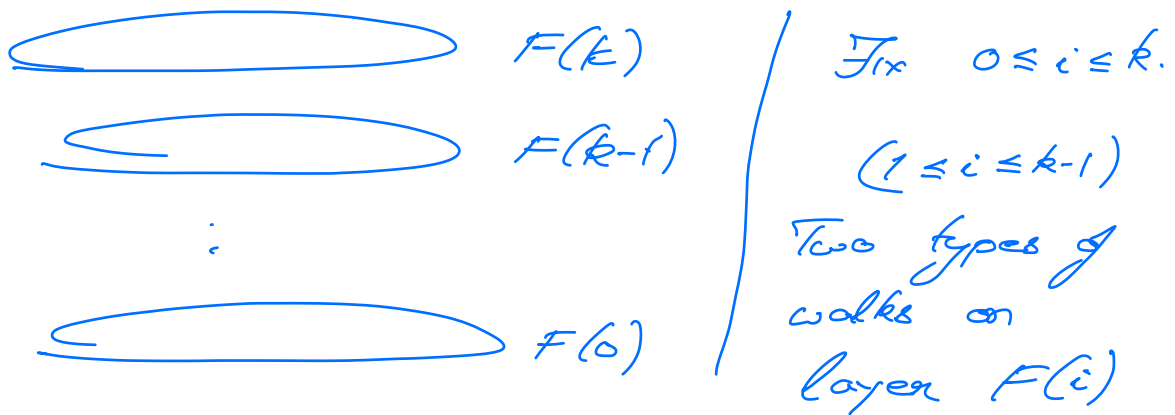




$F = (F(0), F(1), \dots, F(k))$

$\hookrightarrow$  down closed family of sets  
 $F(0) = \{\emptyset\}$

Define: suitable notion of expansion.



Down-Up Walk



Given  $t \in F(i)$   
 - Let  $a \in t$   
 $b \leftarrow t \cdot \{a\}$   
 -  $t' \supseteq b$

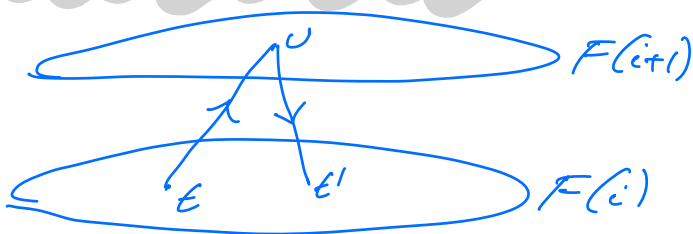
from the dist  $\pi_i / \epsilon'_{\geq i}$

Down-Up Walk on a graph  
on layer  $E(1) = V$



(regular graph  
DU-walk =  $J/n$ )

Up-Down Walk on layer  $i$



Given  $\epsilon \in F(i)$

Pick  $u \leftarrow \pi_{\epsilon+1} / \epsilon \geq \epsilon$

Pick  $\epsilon' \leftarrow \pi_i / \epsilon' \leq u$

Up-down walk on a graph  
on layer  $V = E(i)$   
 $V = E(2)$



Non-Lazy component of the up-down  
walk - usual random  
walk on a graph.

Graph: (regular)

$$\text{---} \quad E(2) = E \quad ($$

$$\text{---} \quad E(1) = V$$

$$\text{---} \quad \cdot \phi \quad E(0) = \{\phi\}$$

$$\|A - J/n\| \leq \lambda$$

$J/n$  - Down-up walk  
on layer  $E(1)$

(ie  $\|A - (J-I)/n-1\| \leq \lambda$ )  $A$  - non-lazy part of  
the up-down walk  
on layer  $E(1)$ .

Generalize this to hypergraph.

Defn:  $(F(k), F(k-1), \dots, F(0))$   
- down closed family of  
sets

is a  $\lambda$ -HDX if

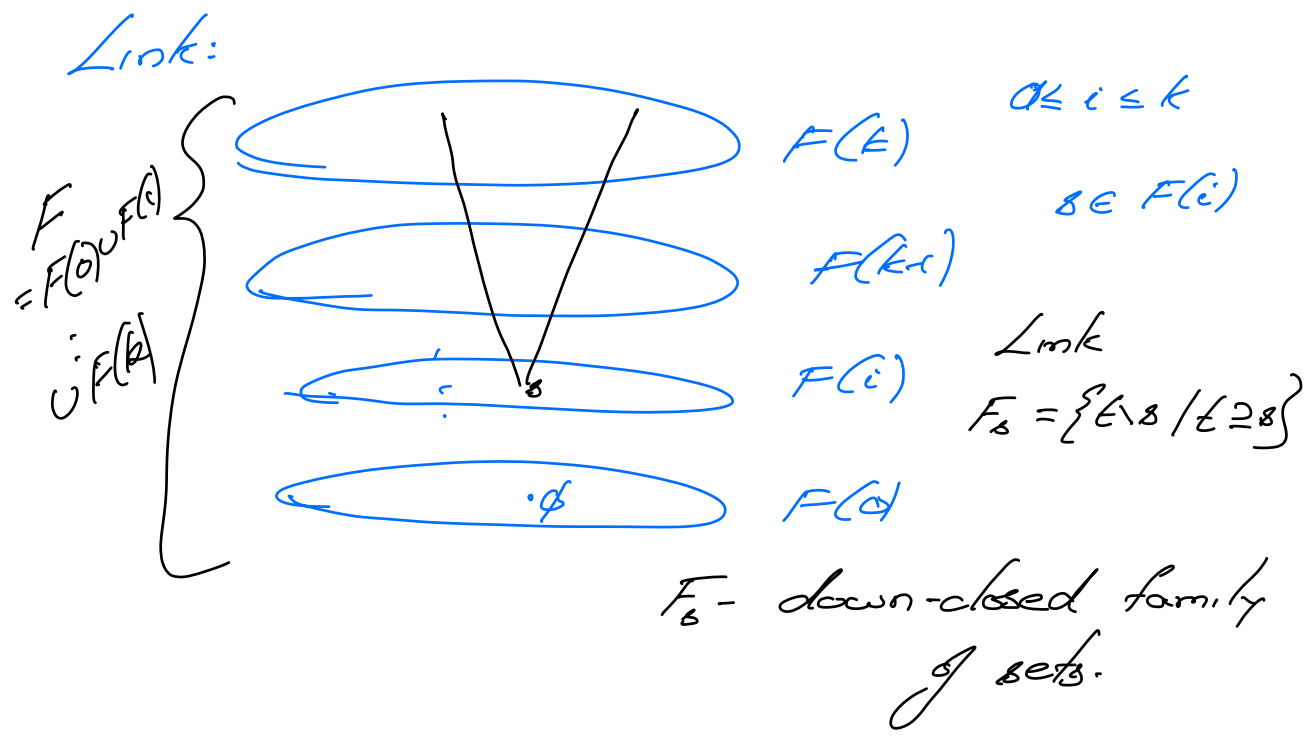
$$\forall 1 \leq i \leq k-1$$

$$\| \text{ri} \cdot P_i^{\wedge} - \text{nl} \cdot P_i^{\vee} \| \leq \lambda$$

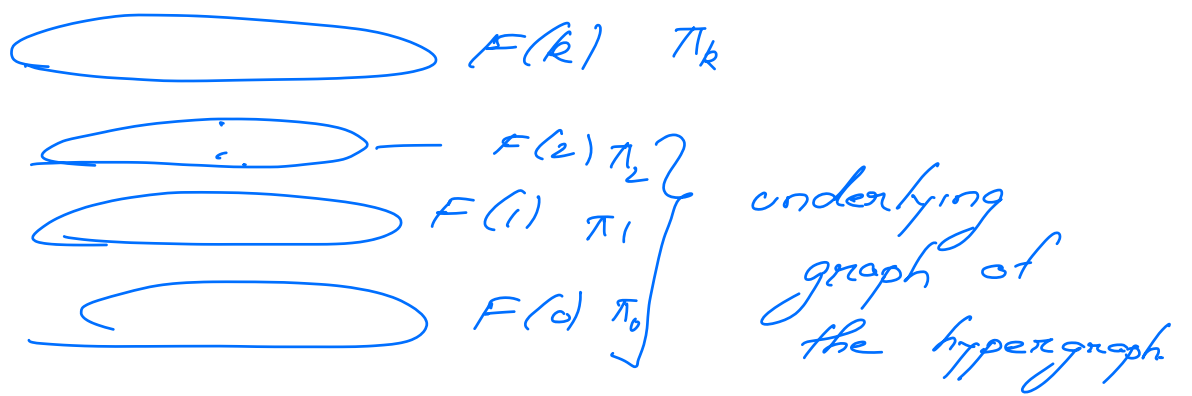
$P_i^{\wedge}$  - up-down walk

$P_i^{\vee}$  - down-up walk

Another defn of HDX  
(in terms of links)



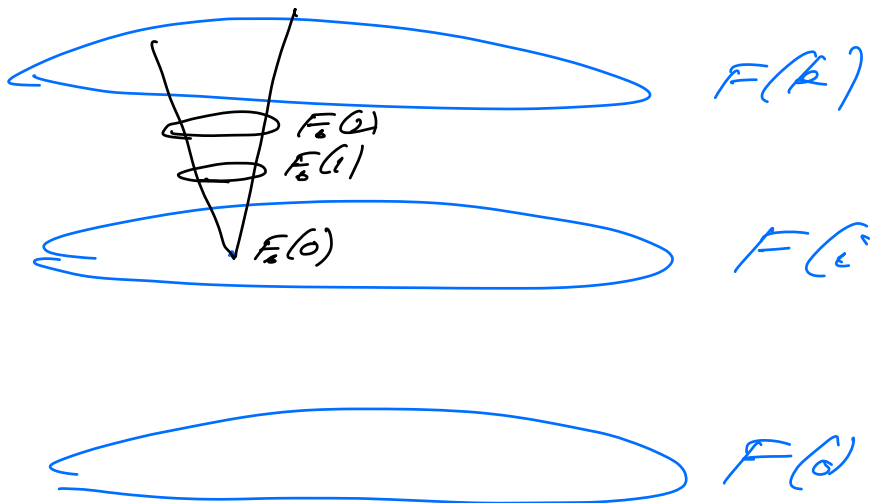
$F = F(0) \cup F(1) \dots \cup F(k).$



Defn.  $F = (F(0), \dots, F(k))$  is a  $\lambda$ -link HDX if [Dionisi-Kaufman]



for all  $0 \leq i \leq k-2$ ,  $\Sigma \in F(i)$   
 the underlying graph of  $F_i$   
 -  $\lambda$ -spectral expander.



Thm:  $\lambda$ -link HDX  $\Rightarrow$   $\lambda$ -HDX  
 conversely for constant  $k$   
 $\lambda$ -HDX  $\Rightarrow$   $k\lambda$ -link HDX

( $k$ -arity of sets in  
 the topmost layer).

Questions:

- (1) Do sparse  $\lambda$ -HDX exist?
- (2) Are HDX "useful"?

(1) Sparse HDX:

- Lubotzky - Samuels - Vishne.  
(Ramanujan Complexes)
- Kautman - Oppenheim (exposition - HS)
- O'Donnell - Pratt

(2) Is this theory of HDX useful?

Up-down / Down-up walks

- help in "counting the number of spanning trees in a graph"

(sampling a random spanning tree).

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Problem:

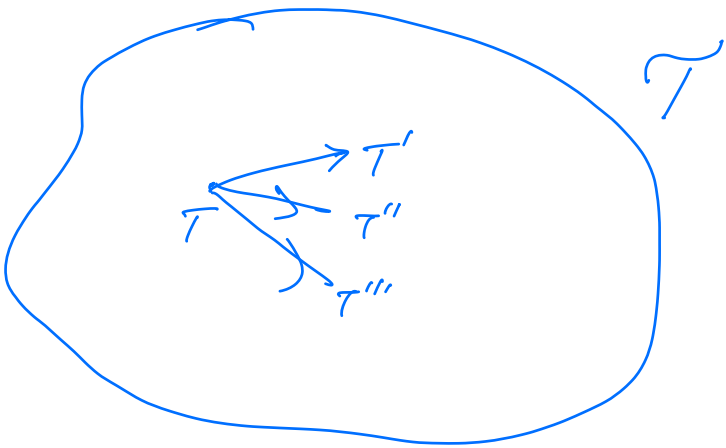
Given a graph (undirected)

$$G = (V, E), \text{ sa}$$

$$\mathcal{T} = \{ T \subseteq E \mid T \text{ is a spanning tree of } G \}$$

sample an elt  $T$  uniformly  $\mathcal{T}$ .

Obs:  $|\mathcal{T}|$  - may be exponentially large in size of graph  $G$ .



Markov chain on  $\mathcal{T}$

Random walk on  $\mathcal{T}$

Random walk converges very fast to the unit dist.

Cluster Dynamics on  $\mathcal{T}$ .

Given a  $T \in \mathcal{T}$ .

- Pick a  $e \in T$

$$S := T \setminus \{e\}$$

- Pick a random  $T' \supseteq S$ .

GD<sub>T</sub>:

$T = \cup \{f\}$   
 for some  $f \in E$   
 s.t.  $\cup \{f\}$  is a  
 tree.

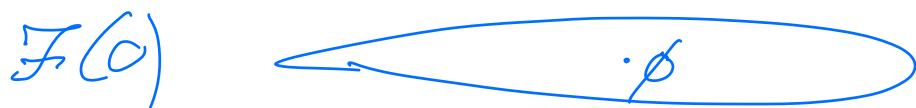
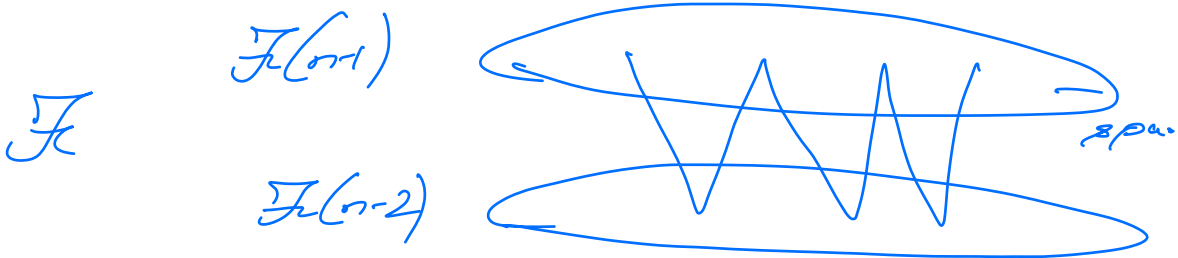
→ Set of all forests in  $G$

$$\mathcal{F} = \{F \subseteq E \mid F \text{ is a forest}\}$$

-  $\mathcal{F}$  - downclosed?

$$\max_{F \in \mathcal{F}} |F| = |V| - 1$$

(assuming  $G$  is connected)



$GD_T$ : Down-Up walk on  $\mathcal{F}$   
on layer  $n-1$ .

To show:  $GD_T$  mixes well

equiv

Down-up walk on layer  $n-1$   
mixes well.

Next lecture: Prove this by showing  
 $\mathcal{F}$  is a good HDX.

[Arani-Liu-Charan-Vinzent '19]